

High speed computer networks

Exterior routing protocols

Sosina M. Addis Ababa institute of technology (AAiT) 2012 E.C.

Routing algorithms

- Distance vector
- Link state
- Protocols
 - RIP and OSPF

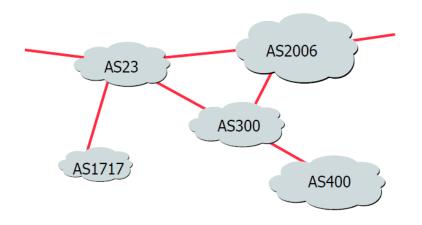
 \checkmark Optimal paths from source to destination that are internal to the same AS

□Not suitable for inter AS routing

• Scalability

- \checkmark LS flooding
- ✓ DV instability

o Metric



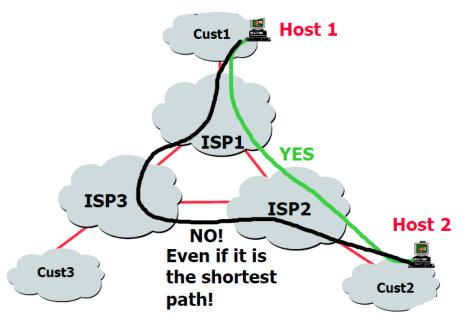
Does not include a distance or cost estimate

Provides path information

• Lists the ASs that must be traversed to reach the destination

• Enables policy routing

✓ E.g., avoid transiting a particular AS



Exterior routing protocol for the internet

Allows gateway routers in different AS to cooperate in the exchange of routing information

Three Functional procedures

 Neighbor acquisition
 Neighbor reachability
 Network reachability

Routers in different ASs wishing to exchange route information

Neighbor acquisition - two neighboring routers in different AS agree to exchange routing information regularly

Procedures

• One router sends a request message to the other

• The other router may either accept or refuse the offer

Given the set of the s

After establishing neighbor reachability

Neighbor reachability - to maintain the relationship

• A periodic keepalive message

Network reachability

- Each router maintains a database of the networks that it can reach and the preferred route for reaching each network
- The router issues a broadcast update message when change is made to the database

Sent over TCP connection

Open message

• The first message after TCP connection establishment

Update message

• Used to transfer routing information between BGP peers

✓ Advertise a feasible route

Keepalive message

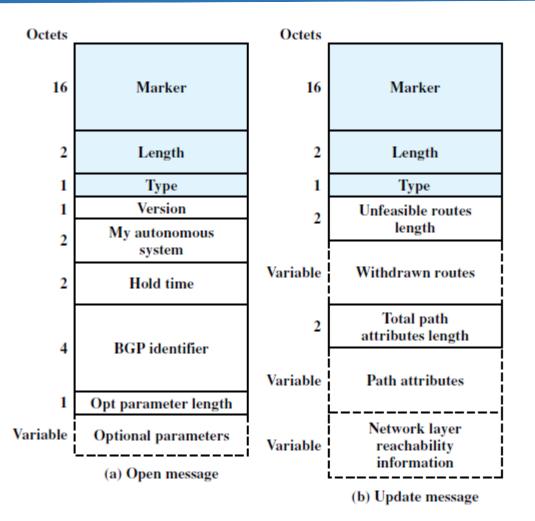
• Sent periodically (every 60 s)

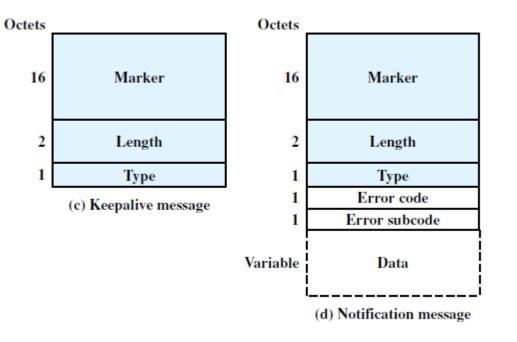
• Wait for a hold-time period (180 s) before declaring the peer is unreachable

□Notification message

• Sent when there is a fatal error condition

BGP messages





- *Marker*: included for compatibility
- *Length*: total length of the message
- *Type*: type code of the message (open, update, notification, keepalive)

A router first opens a TCP connection to a neighbor router

• BGP runs over TCP on a well-know port (179)

Sends an open message

• AS, IP address, Hold time parameter

Recipient

• Calculate hold time (min[it hold timer, hold timer in the open message])

• Reply with keep alive message

 \checkmark consists simply of the header

Update message communicates two types of information

• Information about a single route

- ✓ Network Layer Reachability Information
- ✓ Total Path Attributes Length
- ✓ Path Attributes

• A list of routes previously advertised by this router that are being withdrawn

• May contain one or both types of information

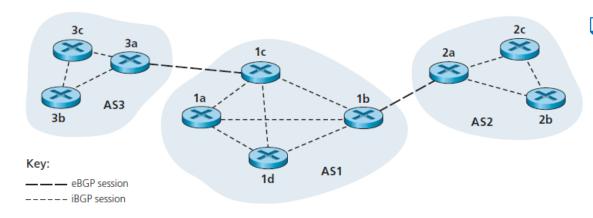
Network Layer Reachability Information(NLRI)

o list of identifiers (IP) of networks that can be reached by this route

Path attributes

- Origin –an interior or exterior router
- AS path Sequence of ASs that are traversed for this route
 - ✓ Prevent looping
 - Enables routing policies implementation
- Next Hop: The IP address of the next-hop router along the path to the destination
- Multi_Exit_disc : Information about routers internal to AS
- Local_Pref : Tell other routers within AS degree of preference for a particular route
- Atomic_Aggregate, Aggregator : Uses subnet addresses in tree view of network to reduce information needed in NLRI

BGP: update



E.g., 138.16.64/24 is first advertised from AS2 to AS1
 AS1 advertises the prefix to AS3: AS-PATH = AS2 AS1

Withdrawal of one or more routes

o route is identified by the IP address of the destination network

Generated when an error condition is detected

- Message header error: authentication and syntax errors
- **Open message error:** errors and options not recognized in an Open or unacceptable Hold Time
- Update message error: syntax and validity errors in an Update message
- Hold timer expired:
- Finite state machine error: any procedural error
- **Cease:** Used by a router to close a connection with another router in the absence of any other error

BGP sessions

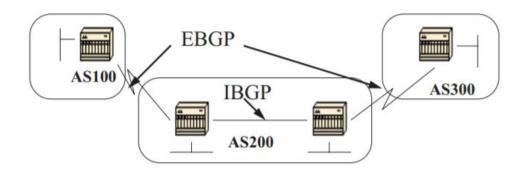
There are two types of BGP sessions

o eBGP: between BGP speaking routers in different AS

✓ Each AS may have more than one router

• iBGP: between BGP routers in the same AS

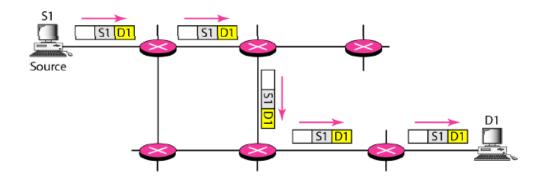
 \checkmark Used to exchange information about external routes



Multicast routing

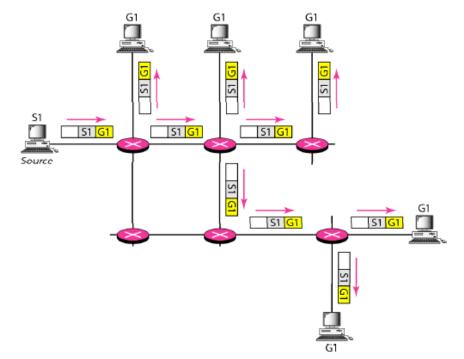
Multicasting versus unicasting

Unicasting



Multicasting

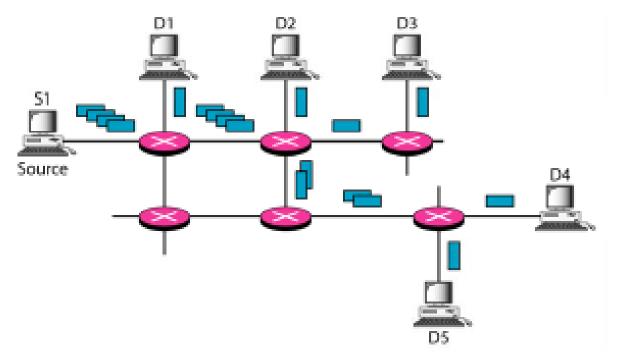
• One source and a group of destination



Multicasting versus multiple unicasting

□ Multiple unicasting

- Several packets start from the source
 - ✓ Not efficient
 - \checkmark May create long delays, particularly with a large group



Access to distributed database

o user's request is multicast to all the database locations, and the location that has the information responds

□Information Dissemination

• For example, a software update can be sent to all purchasers of a particular software package

Dissemination of News

Teleconferencing

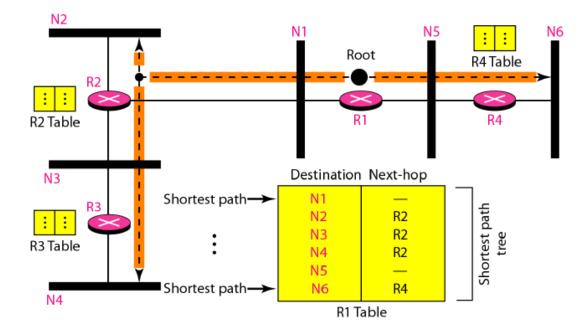
• The individuals attending a teleconference all need to receive the same information at the same time

Optimal routing: shortest path tree

 \odot The root of the tree is the source, and the leaves are the potential destinations

• The path from the root to each destination is the shortest path

unicast routing: the router knows the shortest path for each destination

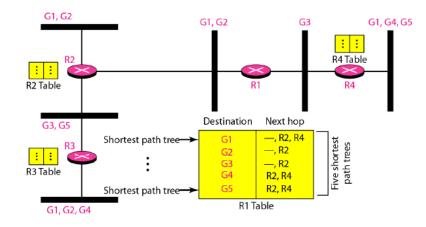


Forwarding of a single packet to members of a group requires a shortest path tree

• If we have *n* groups, we may need *n* shortest path trees

Source-based tree approach

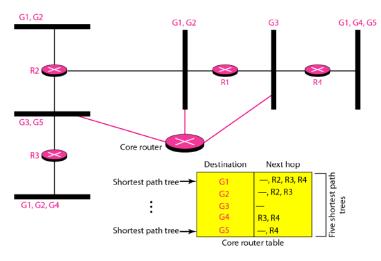
• In the source-based tree approach, each router needs to have one shortest path tree for each group



Complex if we have hundreds or thousands of groups

Group-Shared Tree

- Only one designated (core) router takes the responsibility of distributing multicast traffic
- If a router receives a multicast packet, it encapsulates the packet in a unicast packet and sends it to the core router
- The core router removes the multicast packet from its capsule, and consults its routing table to route the packet



Multicast link state routing: MOSPF

- o uses the source-based tree approach
- Builds a tree that contains all the hosts belonging to a group, the unicast address of the host is used in the calculation
- For efficiency, the router calculates the shortest path trees on demand (when it receives the first multicast packet)

Multicast Distance Vector: DVMRP

uses source-based trees, but the router never actually makes a routing table
a process based on four decision-making strategies

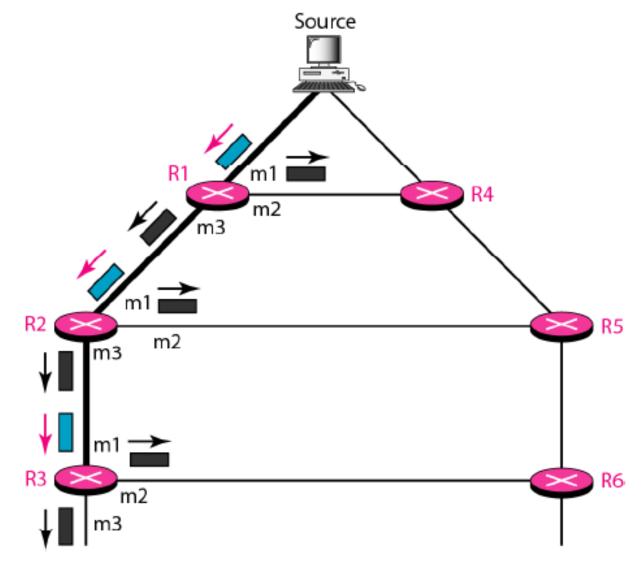
✓ Flooding broadcast packet

- Every network with active members receives the packet.
- However, so will networks without active members
- creates loops in the systems

✓ Reverse Path Forwarding (RPF)

- is a modified flooding strategy
- To prevent loops, only one copy is forwarded; the other copies are dropped
- The router receives a packet and extracts the source address (a unicast address)
- It consults its unicast routing table as though it wants to send a packet to the source address
- The router forwards the packet if it has traveled from the shortest path; it discards it otherwise

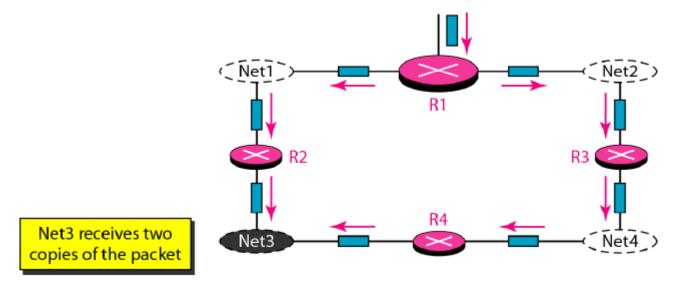
Reverse path forwarding



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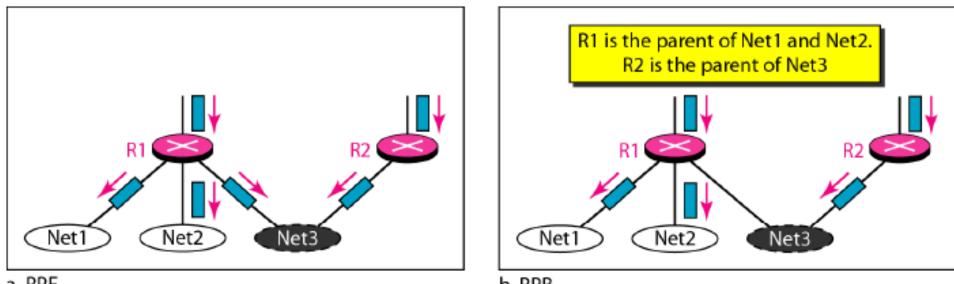
Reverse path broadcasting (RPB)

• RPF does not guarantee that each network receives only one copy



• A network can receive a multicast packet from a particular source only through a designated parent router

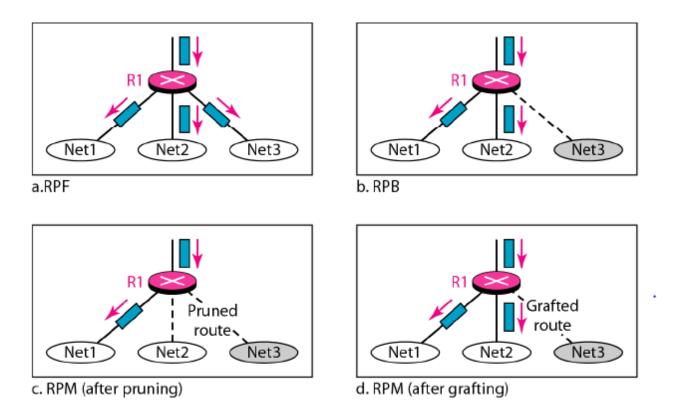
Reverse path broadcasting (RPB)



b. RPB

Reverse path multicasting (RPM)

• To convert broadcasting to multicasting, the protocol uses two procedures, pruning and grafting



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